

Foundations for Actively Secure Card-Based Cryptography

FUN 2020 @ FUN 2022 @ FUNvignana (Not my fav. FUN pun)

Alexander Koch (KIT) and Stefan Walzer (University of Cologne) | 31. May 2022



www.kit.edu

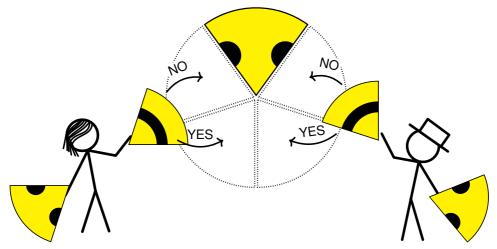


Problems at a Movie Evening...



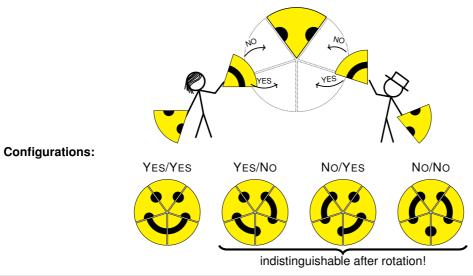
Graphic by S. Walzer, Seats/curtains CC0, Alice/Bob adapted from xkcd (by Randal Munroe) CC BY-NC 2.5, logos copyrighted



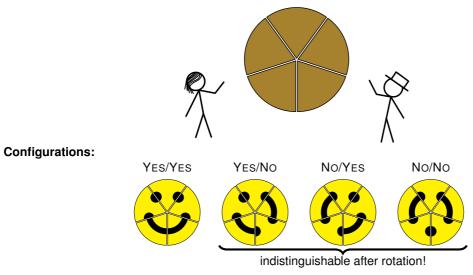


Den Boer (1989). Version with tiles by Verhoeff (2014), Graphic by S. Walzer

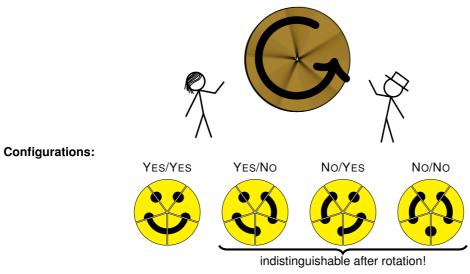






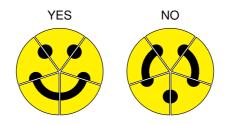






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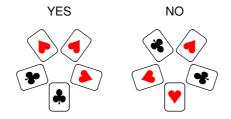
Revealing Tiles...



If a player said NO, then that player cannot know what the other player said!



Revealing Cards... (equivalent)



If a player said **NO**, then that player cannot know what the other player said!



"We have an implementation"





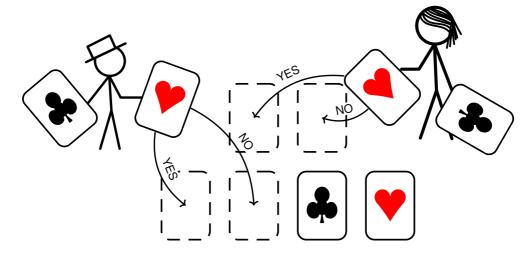
Motivation: Explain Cryptography to Students



By brett jordan via flickr CC BY 2.0



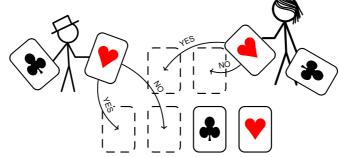
Computing AND with hidden output

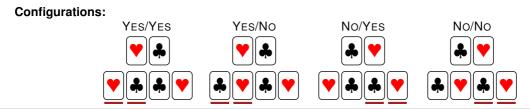


6-Card protocol by Mizuki and Sone (2009)



Computing AND with hidden output

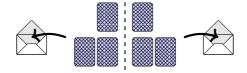




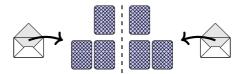
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Computing AND with hidden output



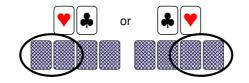




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Computing AND with hidden output





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Research Question

We want to compute arbitrary Boolean circuits

- For this, the output needs to stay hidden, in 2-card encoding
- We need protocols for AND, NOT and Bit-Copy

Question: What are the best protocol with hidden output?

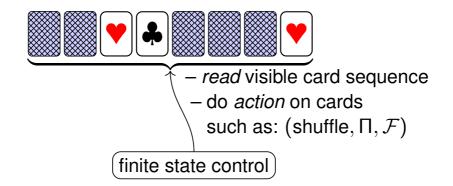
Criteria:

- Number of cards used
- Sunning time behavior (finite vs. Las Vegas)
- Number of steps
- Practically of Shuffling steps

Motivation II: Studying Unconventional Machine Models

For Example, the Model of Mizuki and Shizuya, 2014





Context: Physical Assumptions in Cryptography





Physical Objects





Physical Processes





Main advantages:

- Transparency and ease of understanding for the user
- Strong, otherwise unachievable, security guarantees

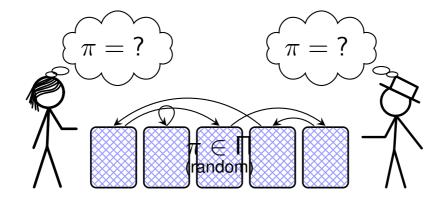
ID card CC0; TAN generator/Scratch-off cards copyrighted; "Miscellaneous Playing Cards" (excerpt) by Philippa Willitts CC BY-NC 2.0; Solar cycle CC0; Schrödinger's cat (excerpt) by ADA&Neagoe CC BY-SA 3.0



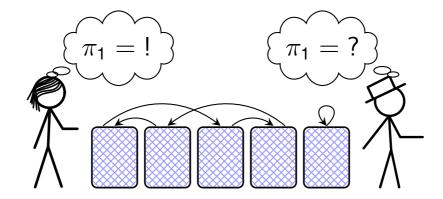
Motivation for our Work

- Mizuki and Shizuya (2014) wanted to capture what can be done with card-based cryptographic protocols, with very general shuffles, but without a description on how they can be performed.
- There have been many ad-hoc solutions, but you have to be careful
- So-far, authors have mainly considered honest-but-curious security
- Our Goal: Implement a large class of shuffles in an actively secure fashion

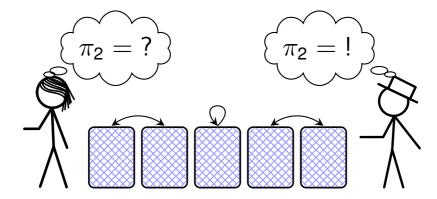




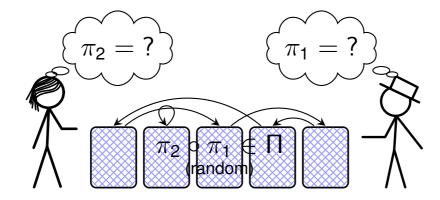














Contribution

Attacker Model

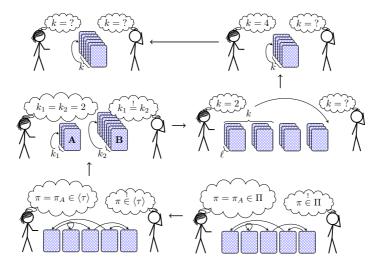
Could maybe be characterized as "Dishonest but slow"

Main Theorem (informal)

Any uniform closed shuffle (i.e. a shuffle with a subgroup as shuffle set, and uniform distribution) can be implemented in the presence of an *active attacker*.

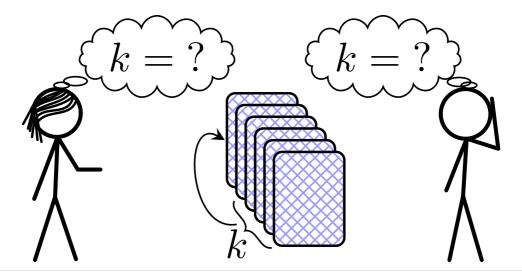


Overview of our Procedure

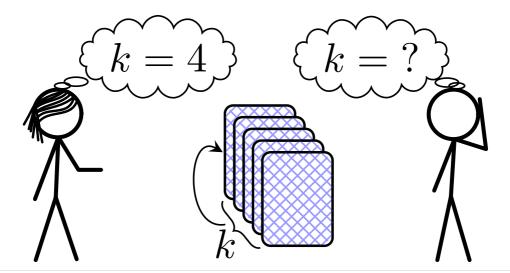




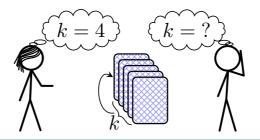
Uniform Cut





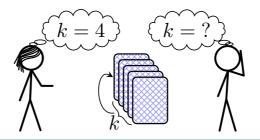






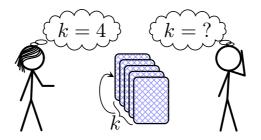


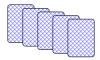




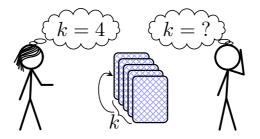


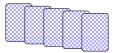




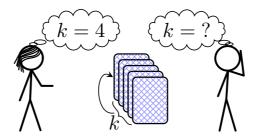






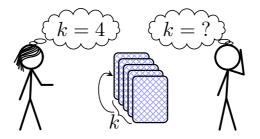


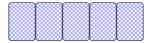




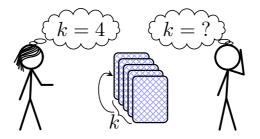


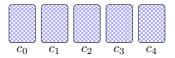




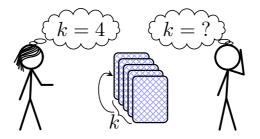


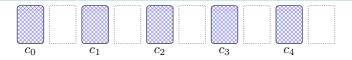




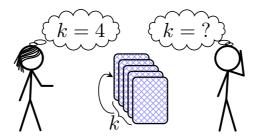






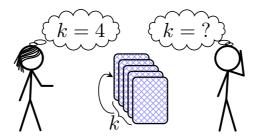


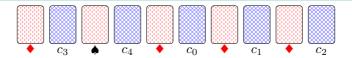




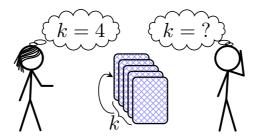


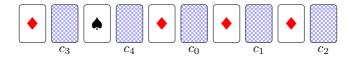




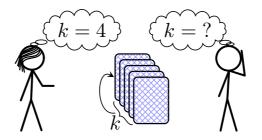


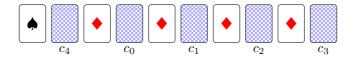




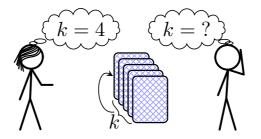


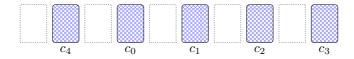




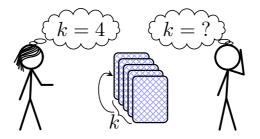


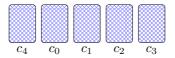




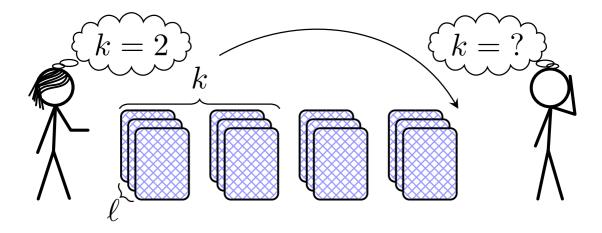




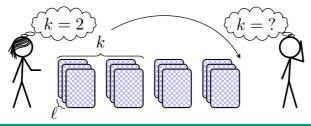




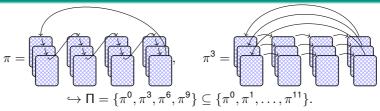




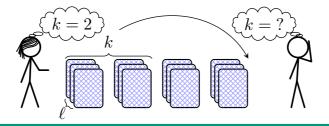


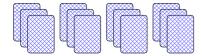


Implementation as Restricted Chosen Cut

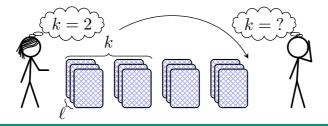






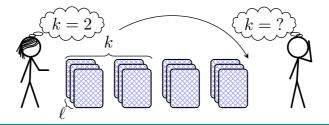


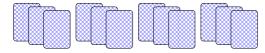




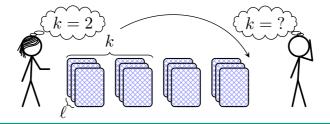


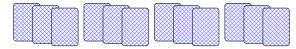




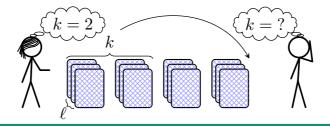


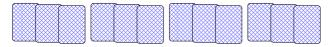




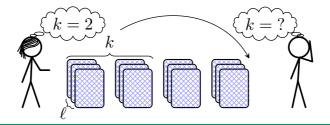






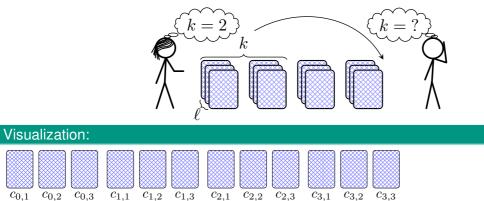




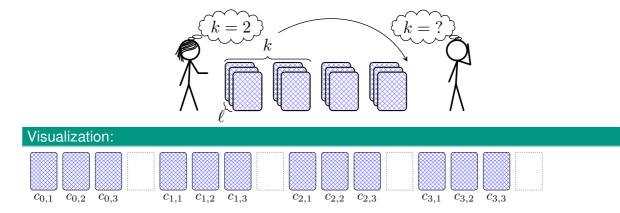




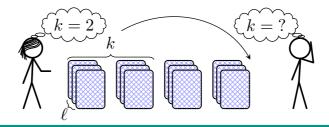


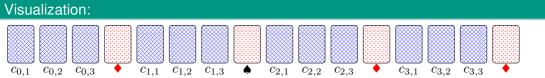




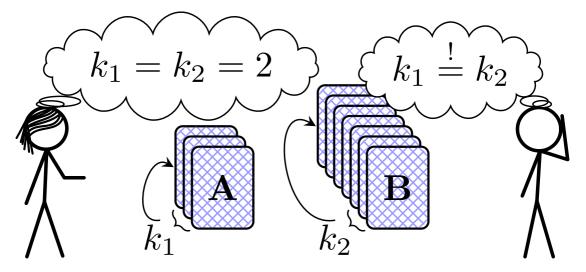




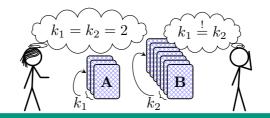


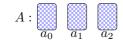


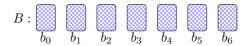




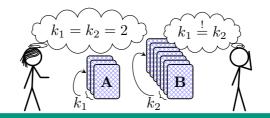


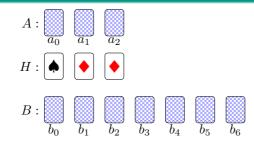




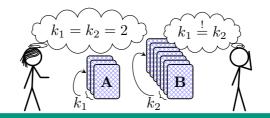


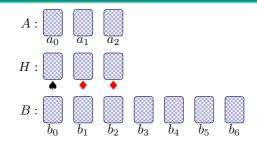




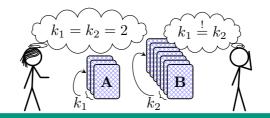


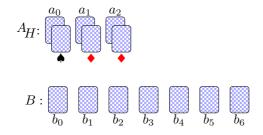




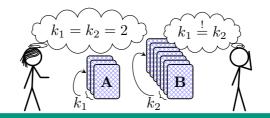


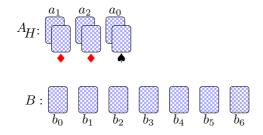




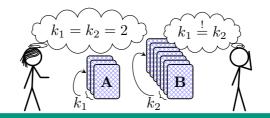


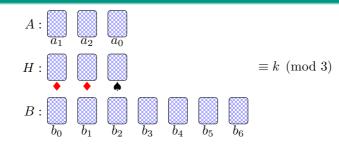




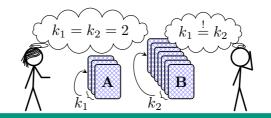


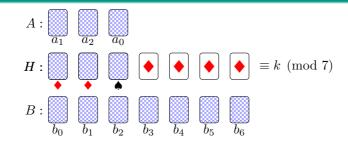




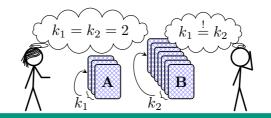


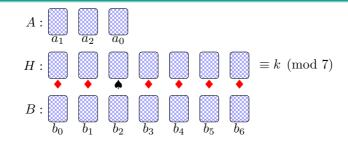




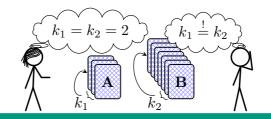


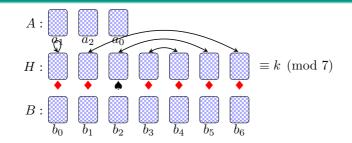




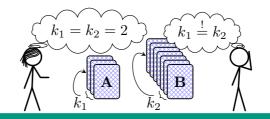


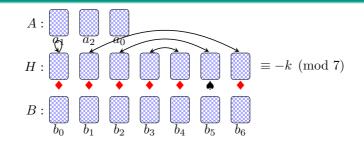




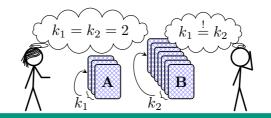


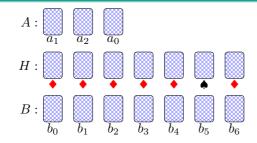




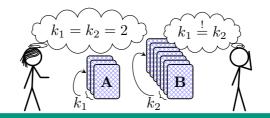


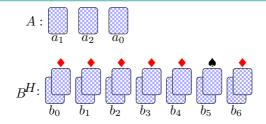




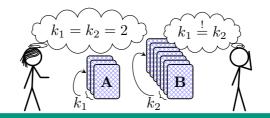


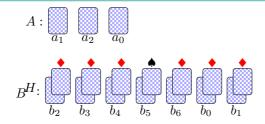




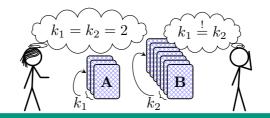


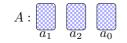


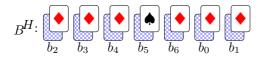




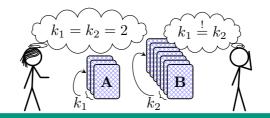


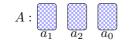


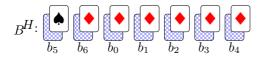






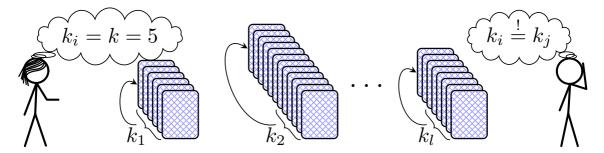








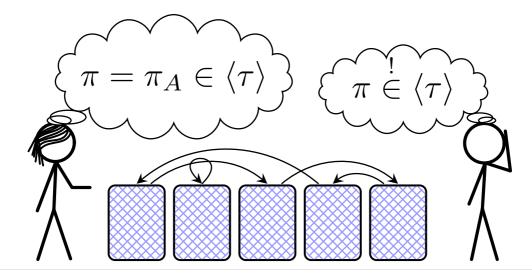
Generalized Coupled Rotation



- Deal with piles one by one.
- Reuse the recorded *k* several times.

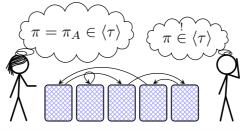


Chosen Permutation from Cyclic Group $\langle \tau \rangle$ for given $\tau \in S_n$





Chosen Permutation from Cyclic Group $\langle \tau \rangle$ for given $\tau \in S_n$



Implementation:

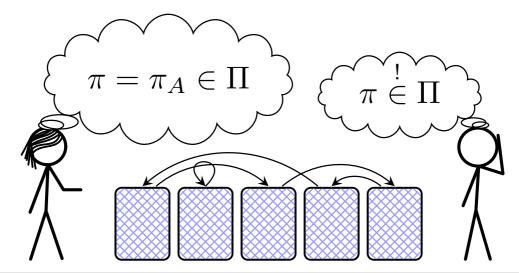
Assume $\tau = (123)(456789)$.

Let $\ell = 3$ be the length of a shortest non-trivial cycle in τ .

- $R = \{\tau^0, \tau^1, \dots, \tau^{\ell-1}\}$ is a generalised coupled rotation.
- $\langle \tau \rangle = \underbrace{R \circ R \circ \ldots \circ R}_{\operatorname{ord}(\tau)/\ell \text{ times}}$
- Perform $\operatorname{ord}(\tau)/\ell$ coupled rotations with *R*.

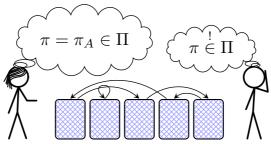


Chosen Permutation from Group





Chosen Permutation from Group



Implementation:

Note that every group can be written as a product of cyclic groups:

$$\Pi = \prod_{\pi \in \Pi} \langle \pi \rangle$$

Now, we can use the previous step. (Admittedly this is not efficient.)



Conclusion and Thanks!

Main take-away

Next time you shuffle in a game (fancier than full S_n), consider active security.

Not covered:

Protocols with permutations encoding the players' inputs and their representation as certain kinds of "state trees" (where security can be verified by a local check at branching nodes and leafs).

Open Problems

- What are all the shuffles that have an actively-secure implementation
- What are all the shuffles that have a *nice and efficient* actively-secure implementation
- Impossibility results for protocols with permutation-encoded inputs.



Literature

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