

BABL

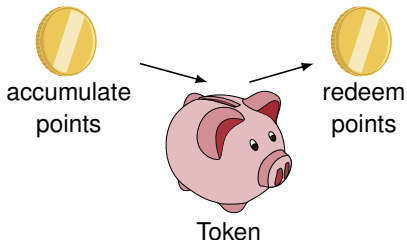
Black-Box Accumulation Based on Lattices

S.H. Faller, P. Baumer, M. Klooß, A. Koch, A. Ottenhues, M. Raiber | 14. Dec 2021



What is Black-Box Accumulation (BBA)?

- Cryptographic protocol for point transfer
- Users can charge their token with points
- Points can be redeemed later
- Introduced by Jager and Rupp [JR16]. Several improvements: [Har+17; Hof+20; Bob+20; Blo+19].



BBA Scenarios

Possible scenarios:

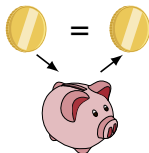
- Prepayment systems (e.g. in public transport)
- Loyalty programs
- Reputation systems

Important BBA Properties

Privacy:



Balance-binding:



Double-Spending-Detection:



Core idea: Updating tokens

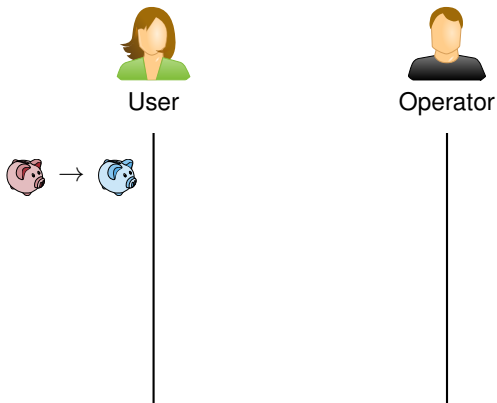


Figure: Sketch of the *Update*-protocol, graphics by M. Nagel

Core idea: Updating tokens

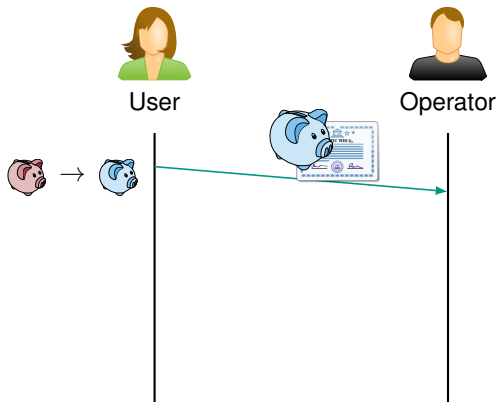


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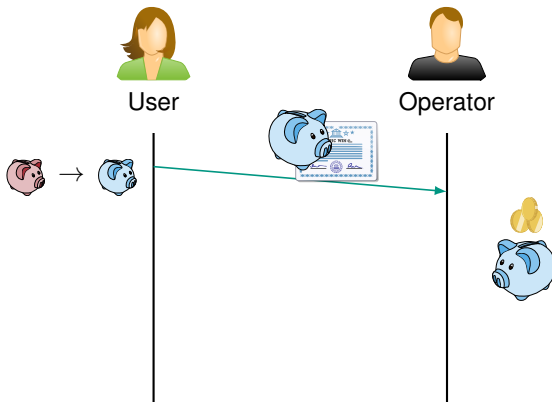


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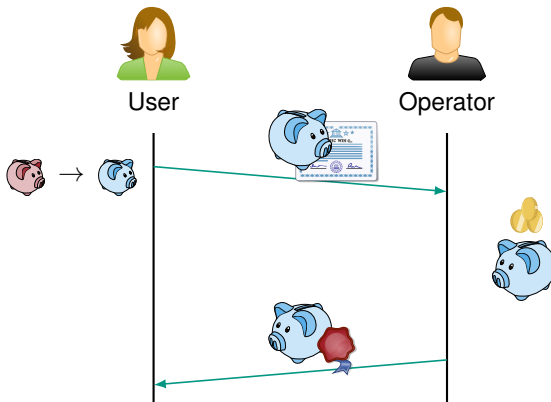


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Motivation

- All prior work on BBA relied on classical hardness assumptions
- Is a lattice-based construction of BBA feasible?
- How efficient will the construction be?

Double-Spending Tags

Same idea as in [Har+17]. Informally:

$$t = \text{sk}_U \cdot u_O + u_U,$$

u_U was fixed when token is issued

u_O is drawn when token is used.

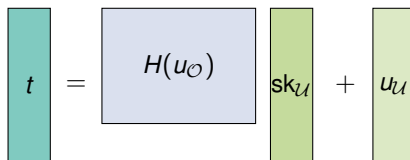
If same token is used again:

$$t' = \text{sk}_U \cdot u'_O + u_U$$

Operators can solve for sk_U .

Store these tags (with a serial number) in a database

Double-Spending Tags



The diagram illustrates the equation for Double-Spending Tags. It consists of a teal vertical rectangle on the left containing the variable t . This is followed by an equals sign. To the right of the equals sign is a light blue square containing the function $H(u_0)$. This is followed by a light green vertical rectangle containing the variable sk_U . To the right of this rectangle is a plus sign, followed by another light green vertical rectangle containing the variable u_U .

$$t = H(u_0) sk_U + u_U$$

SIS Commitments [KTX08]

Let $m = (m_1, \dots, m_k) \in (\{0, 1\}^m)^k$, $D_0, D_1, \dots, D_k \in \mathbb{Z}_q^{n \times m}$, $r \leftarrow \mathbb{Z}_q^m$

$$\text{Com}(m; r) := \begin{array}{|c|} \hline D_0 \\ \hline \end{array} \begin{array}{|c|} \hline r \\ \hline \end{array} + \begin{array}{|c|} \hline D_1 \\ \hline \end{array} \begin{array}{|c|} \hline m_1 \\ \hline \end{array} + \dots + \begin{array}{|c|} \hline D_k \\ \hline \end{array} \begin{array}{|c|} \hline m_k \\ \hline \end{array}$$

Block Substitution

Let $m_i \neq \tilde{m}_i$.

$$c := D_0 \cdot r + D_1 \cdot m_1 + \dots + D_k \cdot m_k$$

$$\text{Subst}(c, m_i, \tilde{m}_i) := c - D_i \cdot m_i + D_i \cdot \tilde{m}_i$$

$\text{Subst}(c, m_i, \hat{m}_i)$ is a valid commitment on \tilde{m}_i

Lattice-based Construction

Example: *Accumulation-Protocol* to get points:

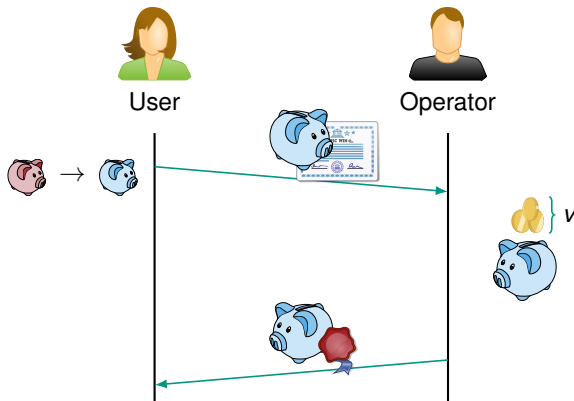


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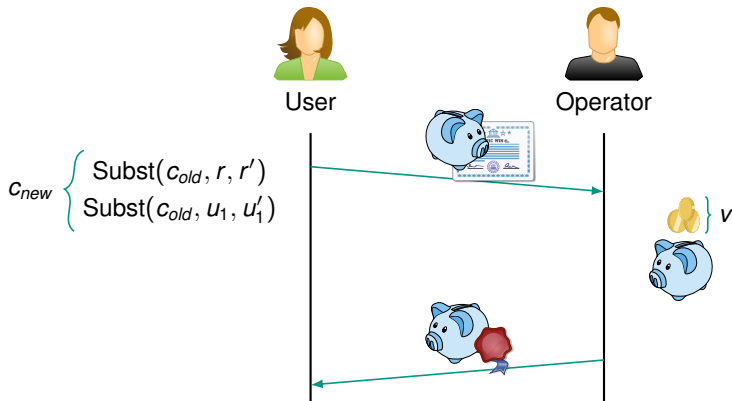


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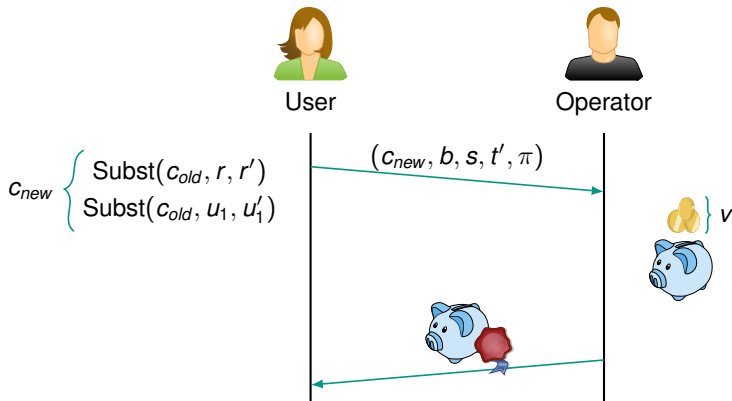


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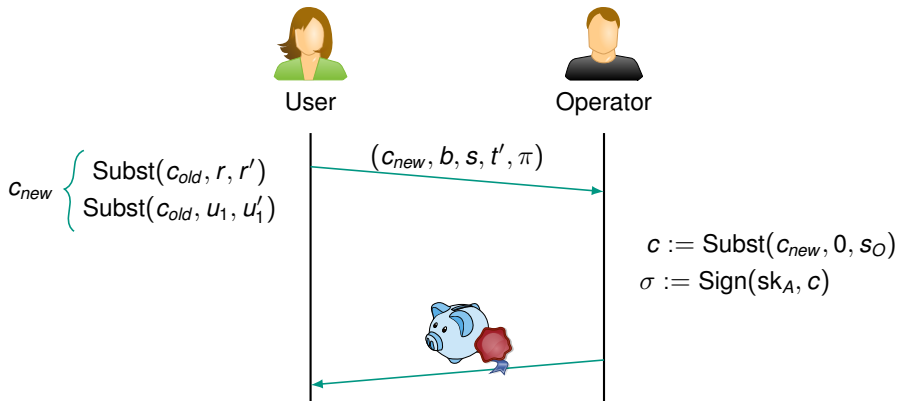


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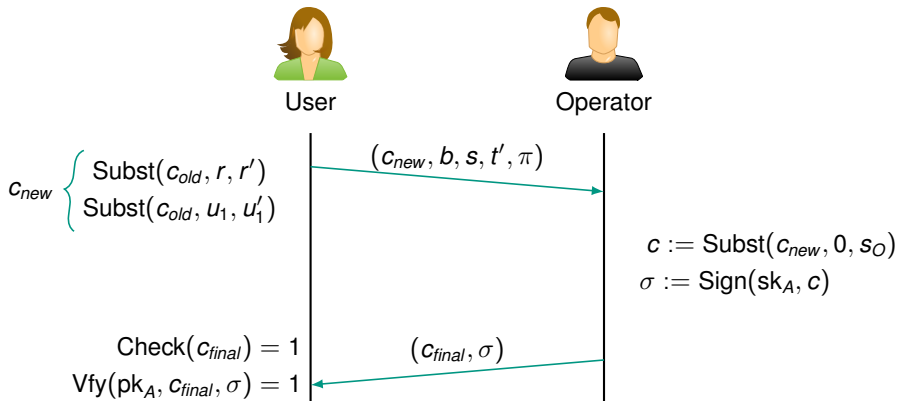
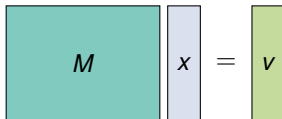


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Zero-Knowledge Proof [Yan+19]

Statement of the form



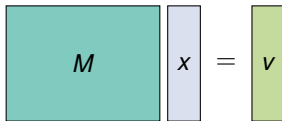
The diagram illustrates the statement of the form $Mx = v$. It consists of three main components: a teal rectangular box on the left containing the letter M , a light blue vertical rectangular box in the middle containing the letter x , and a light green vertical rectangular box on the right containing the letter v . An equals sign is placed between the x box and the v box.

and additionally multiplicative relations:

$$\forall (h, i, j) \in \mathcal{M} : x[h] = x[i] \cdot x[j]$$

Zero-Knowledge Proof [Yan+19]

Statement of the form



The diagram illustrates the statement of the form $Mx = v$. It consists of three rectangular boxes arranged horizontally. The first box is teal and contains the letter M . The second box is light blue and contains the letter x . The third box is light green and contains the letter v . An equals sign is placed between the second and third boxes.

and additionally multiplicative relations:

$$\forall (h, i, j) \in \mathcal{M} : x[h] = x[i] \cdot x[j]$$

E.g. ensure that vectors are binary: $\mathcal{M} = \{(i, i, i) \mid i \in \{1, \dots, n\}\}$.

Zero-Knowledge Proof [Yan+19]

Zero-Knowledge Proofs ensures that:

- Valid signature on token
- Double-spending tag correct
- Commitment contains user's secret key and balance
- Balance is right
- Witness has small norm

Choice of parameters

Used the root Hermite factor estimation formulas of [Yan+19] for SIS/LWE instances

Goal:

$RHF \leq 1.0048$. Corresponds to 80-bit security.

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Instances:

- SIS for binding commitment
- SIS for the public-private key pair
- SIS for the signature to be secure.
- SIS and LWE for the zero-knowledge protocol [Yan+19].

Efficiency estimation

Protocol	Issuance	Transaction	Token	Based on
Our work	70 MB	199 MB	11 MB	Lattices
E-Cash [Lib+17]	33 TB	720 TB	4 MB	Lattices
E-Cash [Yan+19]	53 MB	262 MB	4 MB	Lattices
BBA+ [JR16]	1 kB	14 kB	<1 kB	Elliptic Curves
BBW [Har+17]	1 kB	5 kB	<1 kB	Elliptic Curves

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However: Efficiency of transaction in our construction without ZK 1.8 MB

Results

- ✓ Lattice-based BBA construction
 - Construction is feasible
 - Communication cost high
 - Costs are comparable to other constructions

Future Work

Enhance efficiency

- Use structured lattices
- Find more efficient building blocks (mainly ZK)

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Bibliography I

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